

Performance of RAN Caching in LTE Mobile Networks: a Real Traffic Analysis

Presenter: Tao Lin

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- 1 Introduction
- 2 Traffic Characteristics
- 3 Performance Analysis
- 4 Implications on Real Systems
- 5 Conclusion

Introduction

- What is **RAN Caching**?
 - Content caching on the devices of RAN (Radio Access Networks), e.g., LTE e-NodeB

- What is done in this paper?
 - the **first work** that thoroughly analyzes the performance of **RAN caching** in LTE mobile networks **using real traffic**

Dataset Description

- Over a **7-day** measurement period
- Collected from **3 e-NodeBs** in a commercial LTE network of **China Mobile**
- More than **3200** access LTE users per day
- More than **60 Million** Http sessions

Data Duration	2015.4.10 – 2015.4.16
Number of e-NodeBs	3
Average Users Number per Day	3249
HTTP Requests	62,104,921
Cacheable Http Requests	12,973,637
Total Download Traffic	1776 GB
Total Cacheable Download Traffic	1044 GB

Dataset Description

Field Meaning	Field Meaning
Timestamp	The time when a HTTP request is initiated
User ID	Anonymous user identifier
Content ID	The unique resource locator (URL) of one requested content
HTTP Method	A token obtained from HTTP request message, e.g., GET, HEAD, POST
Content Type	A value obtained from the content type field in HTTP reply message, including text, image, video, audio, application and others
Content Length	The length of requested content in bytes, obtained from HTTP reply message
Downlink Length	Actual length of downlink transmission data
HTTP Status Code	Indication of the implementation state of a specific HTTP session, e.g., 200, 203, 300

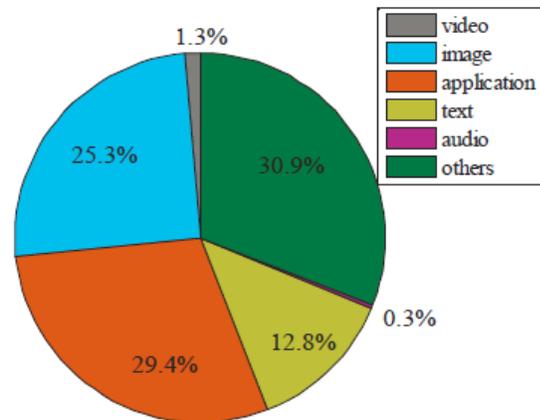
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Main Observations

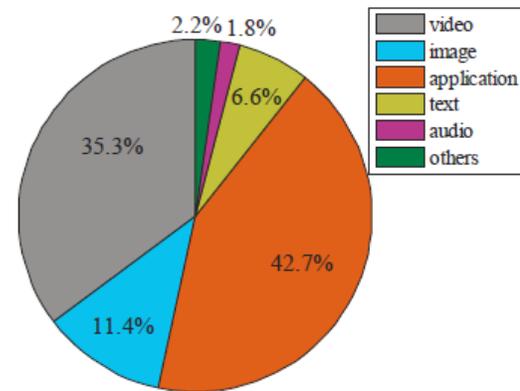
- **Observation 1:** 21.5% HTTP sessions are cacheable. However, the traffic volume of these cacheable HTTP sessions account for 56.2% of the total HTTP traffic

Main Observations

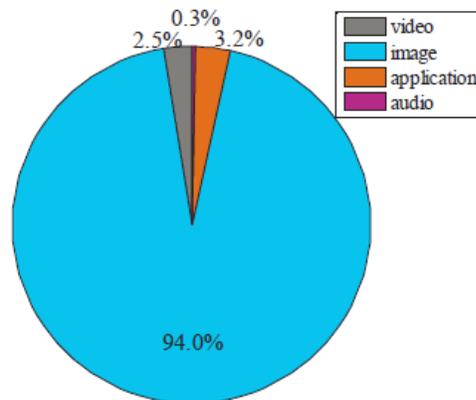
- **Observation 2:** Proportion of Requests and Traffic Volume by Content Type



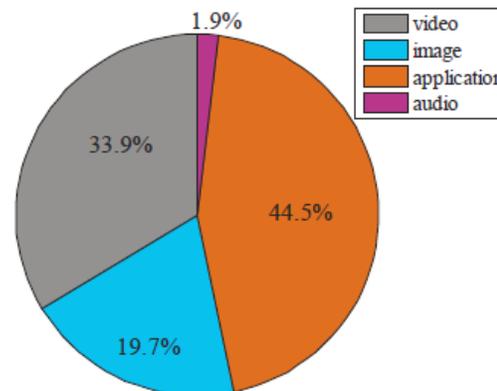
Proportion of All Http Requests



Proportion of All Traffic Volume



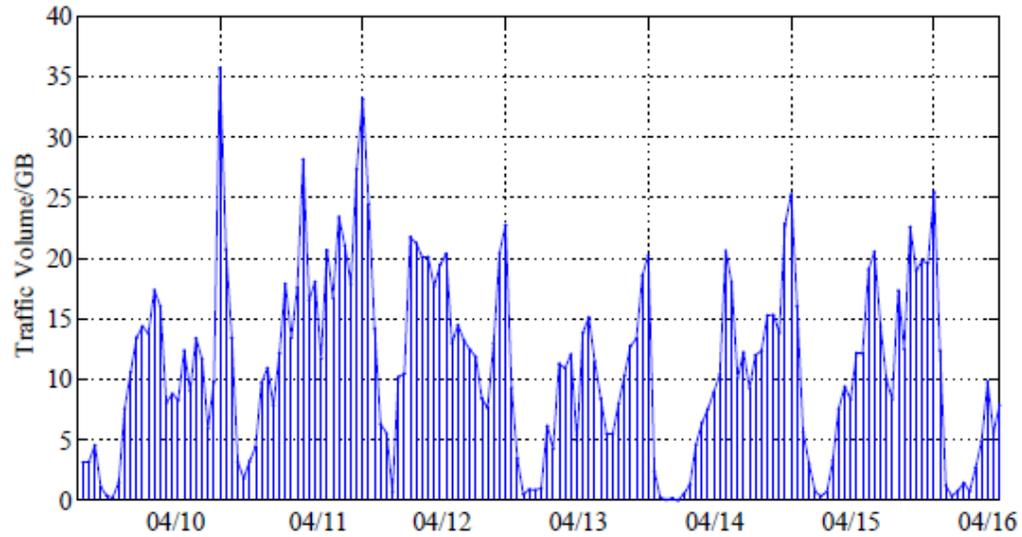
Proportion of Cacheable Http Requests



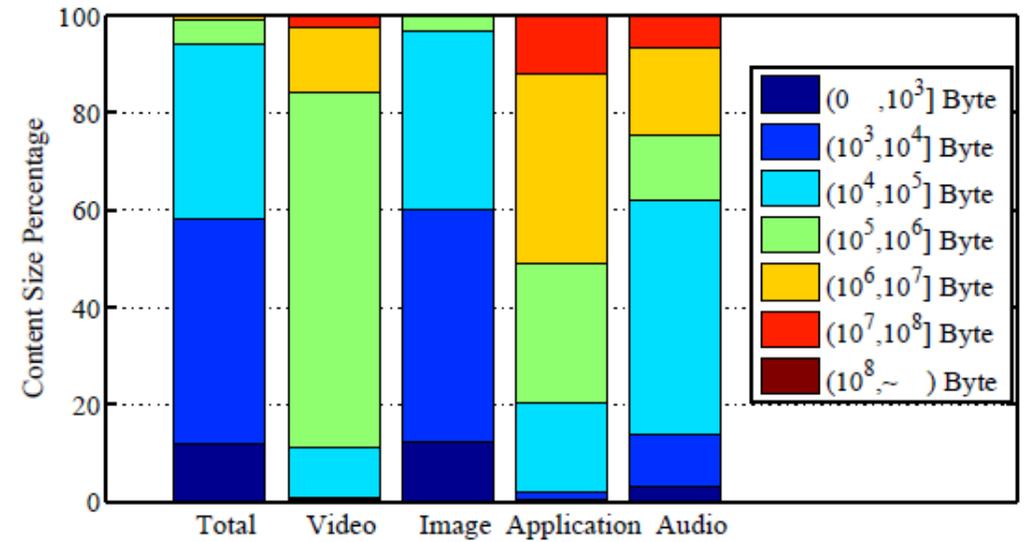
Proportion of Cacheable Traffic Volume

Main Observations

- **Observation 3:** Traffic Volume and Content Size



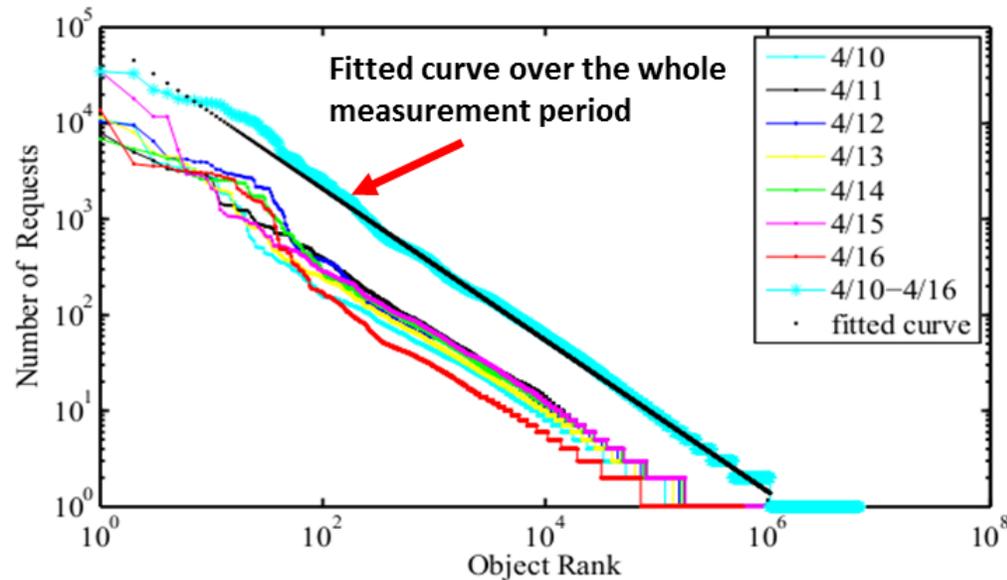
Traffic volume for one week



Cacheable content size distribution

Main Observations

- **Observation 3:** the popularity of cacheable contents follows a **Zipf** distribution with parameter α equal to **0.8**, with 3200 average access users per day.



Cacheable content size distribution

Data	α	C	r
2015/04/10	0.7584	3.8551	0.9747
2015/04/11	0.7425	4.0530	0.9974
2015/04/12	0.8403	4.2719	0.9728
2015/04/13	0.7556	3.9325	0.9892
2015/04/14	0.7885	4.1235	0.9819
2015/04/15	0.7356	3.9844	0.9847
2015/04/16	0.9246	4.1816	0.9568
2015/04/10-16	0.8077	4.9569	0.9925

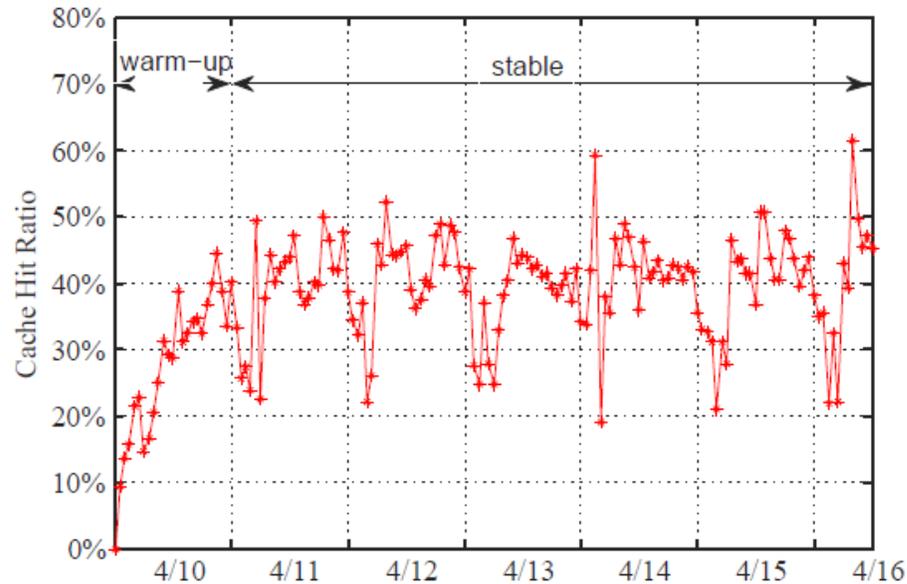
the Obtained Parameters of Zipf Distribution

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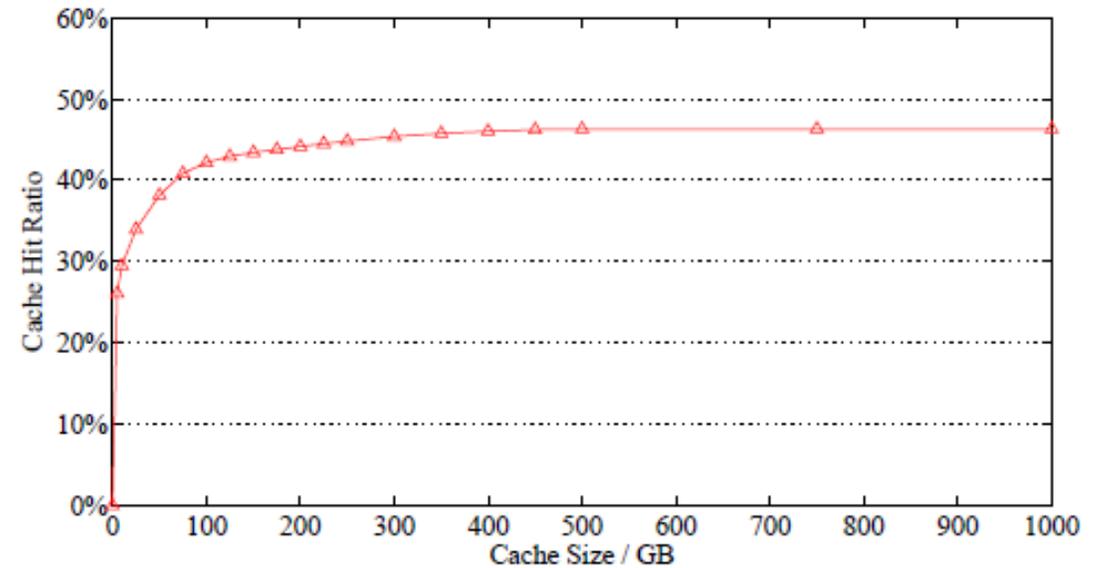
Cache hit ratio based on the total three eNodeBs

- *Effect of Cache size on Hit Ratio*

- ▣ With cache size increasing from 5 GB to 1000 GB, the cache hit ratio rises correspondingly and converges to a maximum hit ratio of **46.3%**.
- ▣ The growth of hit ratio flattens out when the cache size reaches **100 GB**, and the corresponding hit ratio is **42.2%**, which is close to the maximum value.



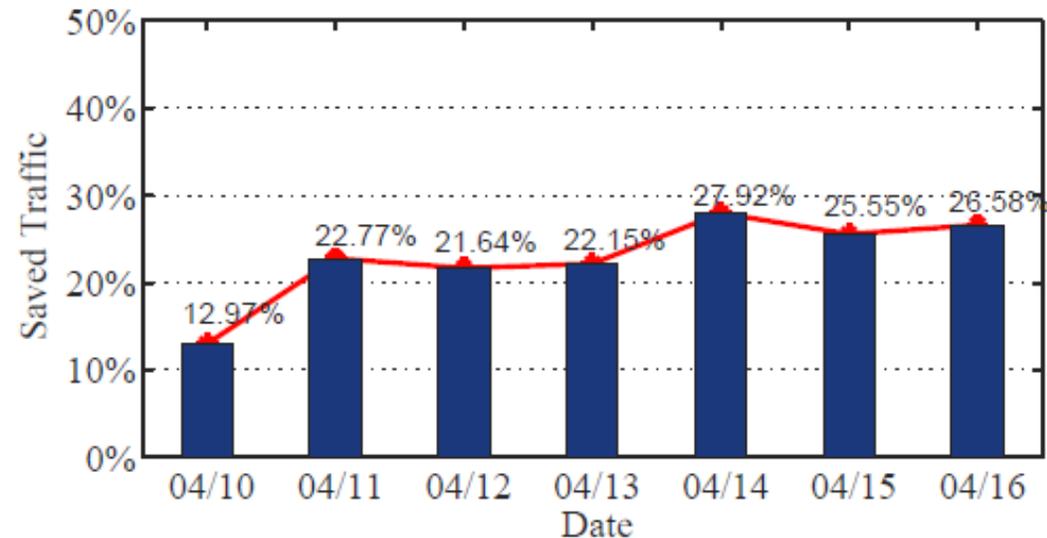
Cache Hit Ratio of the total three eNodeBs



Hit ratio over different cache size

Saved Traffic based on the total three eNodeBs

- *Saved Traffic based on the total three eNodeBs*
 - ▣ With 100 GB cache size, RAN caching can save as much as **24.4%** of the total HTTP downstream traffic (including both cacheable and un-cacheable traffic) on average, with maximum savings of **27.9%**

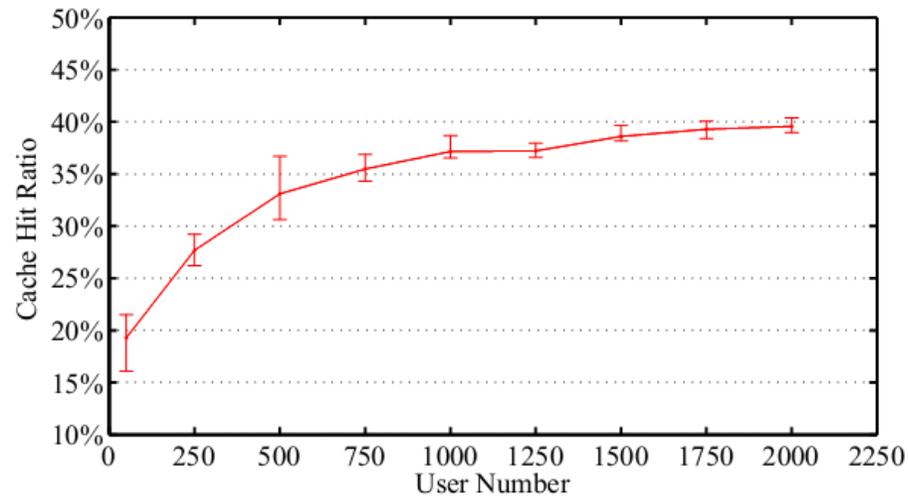


Saved traffic in different days

Effect of User Number

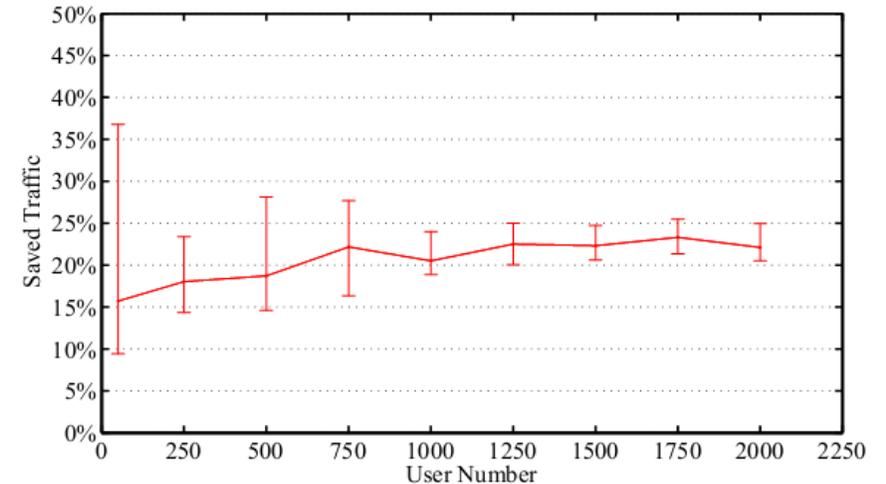
- *Effect of User Number on Hit Ratio*

- ▣ We investigate the effect of user number ranging from 50 to 2000, on the cache hit ratio. Not surprisingly, with increasing user number, the cache hit ratio correspondingly improves from an average value of 19.1% to 39.7%



- *Effect of User Number on Traffic Reduction*

- ▣ The effect of user number on the saved traffic is also investigated. When the user number ranges from 50 to 2000, the volume of total HTTP downstream traffic is reduced by 15.9% to 23.2% on average.



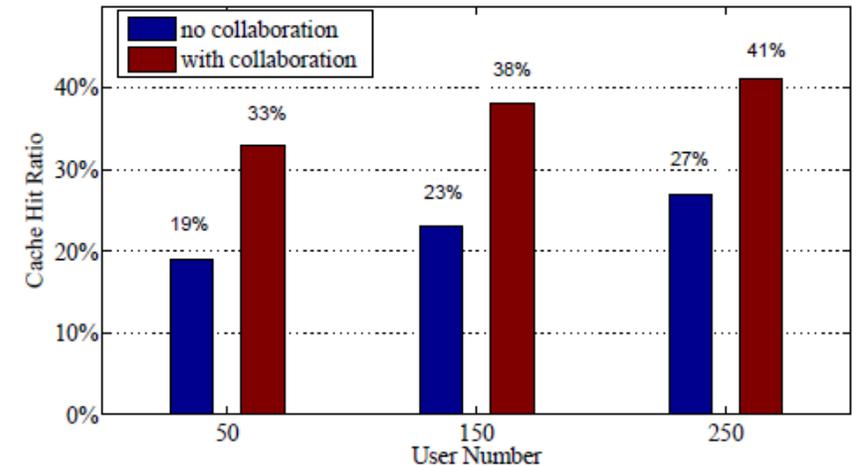
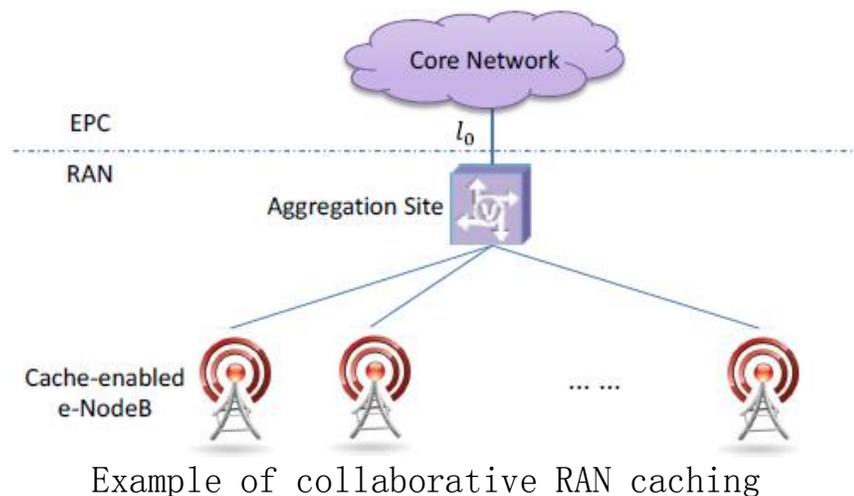
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On Fundamental Design Issues

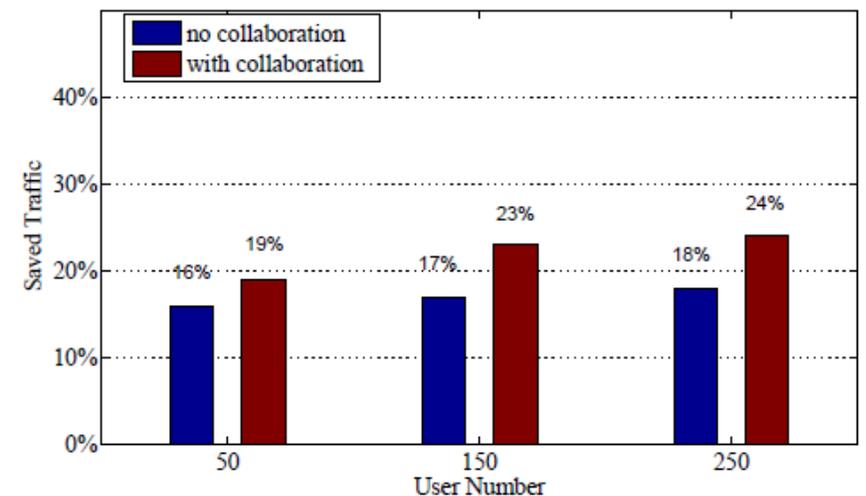
- ❑ What type of content should be stored in the RAN cache?
 - ❑ In addition to **videos** and **applications**, the content type of **image** should not be overlooked, especially considering the vast majority of request proportion (94%) and a remarkable hit ratio (41.9%).
- ❑ How large the storage size of the RAN cache should be?
 - ❑ a cache size of **several hundreds of GB** is enough for a single eNodeB of medium traffic load, and provides a good tradeoff between the performance and the cost of the RAN cache.
- ❑ Which e-NodeB is suitable for deploying RAN caches?
 - ❑ **the access user number** of eNodeB is an important factor when deciding which eNodeBs are good candidates for cache

On the Collaborative Caching Issue

- ❑ What is the potential gain of collaborative RAN caching?
 - ❑ Experiment setup (10 eNodeBs, the user number of each eNodeB set to be 50, 150, 250)
 - ❑ collaborative caching helps improve the performance of a caching system
 - ❑ benefits vs. cost



Comparison of cache hit ratio over different user number



Comparison of saved traffic over different user number

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Conclusion

- the first work that thoroughly analyzes the performance of RAN caching in LTE mobile networks using real traffic
- the popularity of cacheable contents approaches to a **Zipf** distribution with parameter **α equal to 0.8**, with **3200 average access users** per day
- provide evidence of the potential benefits of deploying RAN caches
- the hit ratio of RAN caching reaches **42.2%** on average and that downstream Http traffic can be reduced by **24.4%**, with **100 GB cache size** and around **3200 mobile users**
- present the implications on some fundamental design issues of practical RAN caching systems



Thank You!

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